

University of Illinois at Urbana-Champaign  
Dept. of Electrical and Computer Engineering

## ECE 220: Computer Systems & Programming

### Recursion

## Writing the Fibonacci Sequence

Anyone remember the Fibonacci sequence?

1, 1, 2, 3, 5, 8, 13, ...

Can anyone write the whole sequence?

*(The rest of us can be done for today!)*

How about this way:  $F(0) = 1$   
 $F(1) = 1$   
 $F(N) = F(N - 2) + F(N - 1)$

This answer is a **recursive definition**, a **function defined in terms of itself**.

## Fibonacci Sequence is Well-Defined

Fibonacci:  $F(0) = 1$   
 $F(1) = 1$   
 $F(N) = F(N - 2) + F(N - 1)$

Given this definition, we say that  $F(N)$

- is **well-defined** because
- it **eventually stops recursing** for all  $N \geq 0$ ,
- or, **equivalently**,  $F(N)$  satisfying the equations **is unique** for all  $N \geq 0$ .

## Some Sequences are Not Well-Defined

This sequence is not well-defined:

$G(0) = 1$   
 $G(N) = [ G(N - 1) + G(N + 1) ] / 2$

What can  $G(N)$  be?

1, 1, 1, 1, 1, 1, 1, ...

1, 2, 3, 4, 5, 6, 7, ...

The possibilities are infinite.

**$G(N)$  is not well-defined.**

## Recursive Functions Must Be Well-Defined for Computers

If you

- write a recursive function
- that is not well-defined,
- **don't expect a computer to choose.**

As you know, **computers are dumb.**

Some well-defined recursive functions

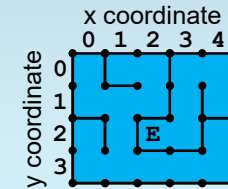
- may still be difficult or impossible to express
- in a computer language.

## Time for the Today's "Help Prof. Lumetta" Problem!

*I need your help again.*

If I build a new house

- as a maze ...
- ... like that ...
- **can I go from (0,0)**
- **to the exit at E?**



## How Did you Solve the Maze?

How did you know?

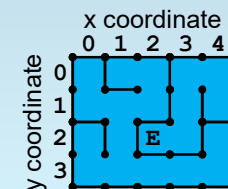
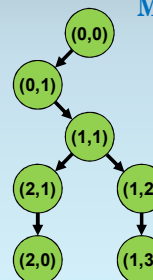
And, more importantly...

...can you teach my computer?

## We Can Use a Tree to Solve the Problem

Maybe you used a tree?

(When we run out of ways to go, we won't have found the exit.)



## Represent the Maze with an Array of Bit Vectors

We can **represent the maze with an array**:

```
static uint8_t maze[10][10];
```

Each space in the array is a bit vector composed of the following bits:

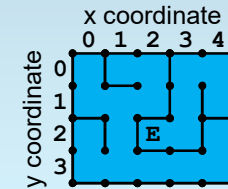
- // 1 – the space has a **left wall**
- // 2 – the space has a **right wall**
- // 4 – the space has an **upper wall**
- // 8 – the space has a **lower wall**
- // 16 – the space is the **exit**

## Do You Understand the Representation?

(Reminder: L=1, R=2, U=4, D=8, E=16)

For example,

- maze[0][0] is 7 (1 | 2 | 4)
- maze[0][1] is 9
- maze[3][1] is 3
- maze[4][2] is 7
- maze[2][2] is 29

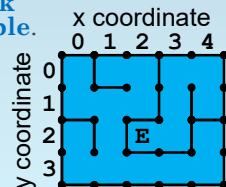


## Outline for a Recursive Solution

Let's solve the problem recursively.

Here's the approach:

- **Keep track of reachable locations.**
- Write a **function to mark one location as reachable.**
- Within the function, **call the same function** to mark all "children" (adjacent reachable neighbors) as reachable.



## Represent Reachable Locations with a Second Array

Track reachable locations with a second array:

```
static uint8_t found[10][10];
```

Each element is either:

- 0 – the space has **not** been **found/reached**
- 1 – the space has been **found/reached**

And we use one variable for the exit:

```
static int32_t saw_exit;
```

(Both of these should be initialized to all 0s.)

## Ready to Write the Recursive Function

Now we're ready to write the function.

Here's a signature:

```
void can_reach (int x, int y);
```

The function should

- set all locations reachable from  $(x, y)$  to 1 in `found`, and
- set `saw_exit` to 1 iff the exit is reachable from  $(x, y)$ .

(To do so, the function will call itself.)

## Mark as Reachable, then Check Children

```
void can_reach (int x, int y)
{
    found[x][y] = 1;
    if (0 == (maze[x][y] & 1)) {
        can_reach (x - 1, y);
    }
    if ((0 == maze[x][y] & 2)) {
        can_reach (x + 1, y);
    }
}
```

$(x, y)$  is reachable.

No left wall?

Space to left is reachable.

## Same Check and Marking for Right Child (Value 2)

```
void can_reach (int x, int y)
{
    found[x][y] = 1;
    if (0 == (maze[x][y] & 1)) {
        can_reach (x - 1, y);
    }
    if (0 == (maze[x][y] & 2)) {
        can_reach (x + 1, y);
    }
}
```

No right wall?

Space to right is reachable.

## Same Check and Marking for Upper Child (Value 4)

```
if (0 == (maze[x][y] & 4)) {
    can_reach (x, y - 1);
}
if (0 == (maze[x][y] & 8)) {
    can_reach (x, y + 1);
}
if (0 != (maze[x][y] & 16)) {
    saw_exit = 1;
}
}
```

No upper wall?

Space above is reachable.

## Same Check and Marking for Lower Child (Value 8)

```

if (0 == (maze[x][y] & 4)) {
    can_reach (x, y - 1);
}
if (0 == (maze[x][y] & 8)) {
    can_reach (x, y + 1);
}
if (0 != (maze[x][y] & 16)) {
    saw_exit = 1;
}
}

```

No lower wall?

Space below is reachable.

## Finally, Check and Mark Exit (Value 16)

```

if (0 == (maze[x][y] & 4)) {
    can_reach (x, y - 1);
}
if (0 == (maze[x][y] & 8)) {
    can_reach (x, y + 1);
}
if (0 != (maze[x][y] & 16)) {
    saw_exit = 1;
}
}

```

Record having seen exit.

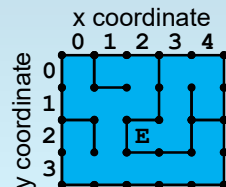
Exit here?

## Does Our `can_reach` Function Answer My Question?

How do we use `can_reach` to answer my question about getting from (0,0) to the exit?

1. Fill `found` and `saw_exit` with 0s.
2. Call `can_reach (0, 0)`.
3. Check `saw_exit`.

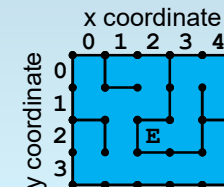
Does it work?



## Maybe There's a Bug?

Let's try it!

(Are you still mad about my asking you to write all of Fibonacci?)



## When Should We Stop Recursing?

### What's the problem?

In math,

- we used base cases
- to stop the recursion.

In a **C** function,

- we **need a stopping condition**
- to stop the recursion.

**So: when should we stop?**

## Stop if the Space Has Already Been Marked as Reachable

Stop if we already reached (x,y).

```
void can_reach (int x, int y)
{
    if (found[x][y]) { return; }
    found[x][y] = 1;
    if (0 == (maze[x][y] & 1)) {
        can_reach (x - 1, y);
    }
}
```

## Review: Induction, Bit-Slicing, Recursion

As mentioned in 120, the following are closely related mathematically

- proof by induction
- bit-sliced hardware design
- recursion.

In all three, one

- solves a small piece of a problem, then
- combines it with the “rest” of the solution
- (which is also solved as small pieces).

## A General Strategy for Recursion

Here's a general strategy for recursion.

```
_____ recursive ( _____ )
{
    // Check stopping conditions.
    // Handle one node.
    // Handle children.
}
```

These may be swapped.