University of Illinois at Urbana-Champaign Dept. of Electrical and Computer Engineering

ECE 120: Introduction to Computing

Boolean Expression Terminology

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Let's Review and Define Some New Terms

literal a variable or its complement examples: A, A', B, B', C, C' sum several terms ORed together

> examples: A + B, AB + B(C + D) + A'C, $A'B' + D(A \oplus B)(C + A')$

product several terms ANDed together

examples: AB, (A + B)(B + CD)(A' + C), $(A' + B')(D + (A \oplus B) + CA')$

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Minterms Were Useful for Proving Logical Completeness

minterm on N inputs

a product in which each variable or its complement appears exactly once (no other factors)

examples: AB', A'B, AB (on inputs A, B)

AB'C, AB'C', A'BC' (on inputs A, B, C)

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A Maxterm Produces a Function with One Zero Row

maxterm on N inputs

a sum in which each variable or its complement appears exactly once (no other terms)

examples: (A + B'), (A' + B), (A + B)

(on inputs **A**, **B**)

(A + B' + C), (A + B' + C'),

(A + 'B + C')

(on inputs A, B, C)

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Sum-of-Products (SOP) Form is Quite Common

sum-of-products (SOP)

a sum (OR) of products (AND) of literals

examples: AB + BC,

AB' + C + A'C'D', but NOT A(B + C) + D

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Product-of-Sums (POS) Form is Also Common

product-of-sums (POS)

a product (AND) of sums (OR) of literals

examples: (A + B)(B + C),

 $(\mathbf{A}+\mathbf{B'})\mathbf{C}(\mathbf{A'}+\mathbf{C'}+\mathbf{D'}),$

but NOT (A + BC)D

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Canonical Forms Allow Easy Comparison, But Are Too Big

canonical SOP

a sum of minterms; the expression produced by the logical completeness construction

canonical POS

a sum of maxterms

What does canonical mean?

Unique (if we assume an ordering on variables).

Too many terms to be of practical value.

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Do You Know Mathematical Implication?

What does A→B mean?

A implies B.

In other words: if A is true, B is also true.

What if **A** is false?

In that case, is $A \rightarrow B$ true or false?

If A is false, $A \rightarrow B$ is true.

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So the Following Odd Statements are True

All purple elephants can fly. (X is a purple elephant \rightarrow X can fly.)

Students who score **above 125**% in ECE120 fail the class.

(X scored **above 125**% → X fails.)

In both, the premise is false for any X, so the implications are true.

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One Function Can Imply Another

A function G is an implicant of a second function F iff G operates on the same variables as F and $G \rightarrow F$.

In other words, every row

- with an output of 1 in **G**'s truth table
- also has an output of 1 in **F**'s truth table.

0 rows in **G**'s truth table do not matter.

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For Our Purposes, Implicants are Products of Literals

In digital design, we only refer to products of literals as implicants.

So we will assume that an implicant can be written as a product of literals.

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We Can Use Implicants to Simplify Functions

As a first step towards simplifying a function **F**, we can ask:

Given an implicant G of F, can we remove any of its literals and obtain another implicant of F?

For example, take F = AB'C + ABC' + ABC.

The first term (AB'C) is an implicant.

Can we remove any literals?

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Try to Remove Each Literal to Find Only AC Implies F

Start from	A	В	\mathbf{C}	F	B'C	AC	AB'
AB'C and try	0	0	0	0	0	0	0
to remove each literal.	0	0	1	0 •	-1	0	0
B'C is not an	0	1	0	0	0	0	0
implicant.	0	1	1	0	0	0	0
AC is an	1	0	0	0 •	0	0	$\overline{(1)}$
implicant.	1	0	1	1	1	1	1
AB' is not an implicant.	1	1	0	1	0	0	0
_	1	1	1	1	0	1	0

We Remove as Many Literals as We Can

So we can simplify **F** by replacing **AB'C** with **AC**:

$$F = AC + ABC' + ABC$$

Checking the second term (ABC'), we find that we can eliminate C' to obtain:

$$F = AC + AB + ABC$$

In the third term (ABC), we can eliminate B or C, but not both. Let's pick B.

$$F = AC + AB + AC$$

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Prime Implicants Have a Minimal Number of Literals

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F = AC + AB + AC

But now we have a duplicate term, which we can eliminate to arrive at a simple form for **F**:

$$F = AC + AB$$

We can remove no more literals.

One more definition: An implicant G of F is a prime implicant of F iff none of the literals in G can be removed to produce other implicants of F.

AB and AC are prime implicants of F.

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To Simplify, Write Function as a Sum of Prime Implicants

The conclusion is obvious:

To simplify a function F, write it as a sum of prime implicants.

Enjoy the algebra.

Good luck!

(Next time, we'll develop a graphical tool that lets us skip the algebra.)

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